



Getting Started on Pink

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Topics

- **Overview: Hardware & Software Architecture**
 - What makes Pink different from other LANL systems.
 - Why we are making it different.
- **Porting Considerations**
- **Compiling**
- **How to Run Jobs**
- **Filesystems**

Aviso: all of this is new (to me, too) and evolving.

- **Lecture assumes familiarity with LANL computing environment.**



Pink Has Two Purposes

- Open production system for Institutional Computing workloads.
- A gateway to future computing strategies at LANL, supporting the following objectives:
 - Better price-performance platforms: Intel Xeon & Myrinet
 - Production-quality, open-source software: RH Linux
 - Vendor-independent HPC features: Linux NetworX, Panasas, CCS-1
 - High-availability cluster computing environment: ClusterMatic Science Appliance

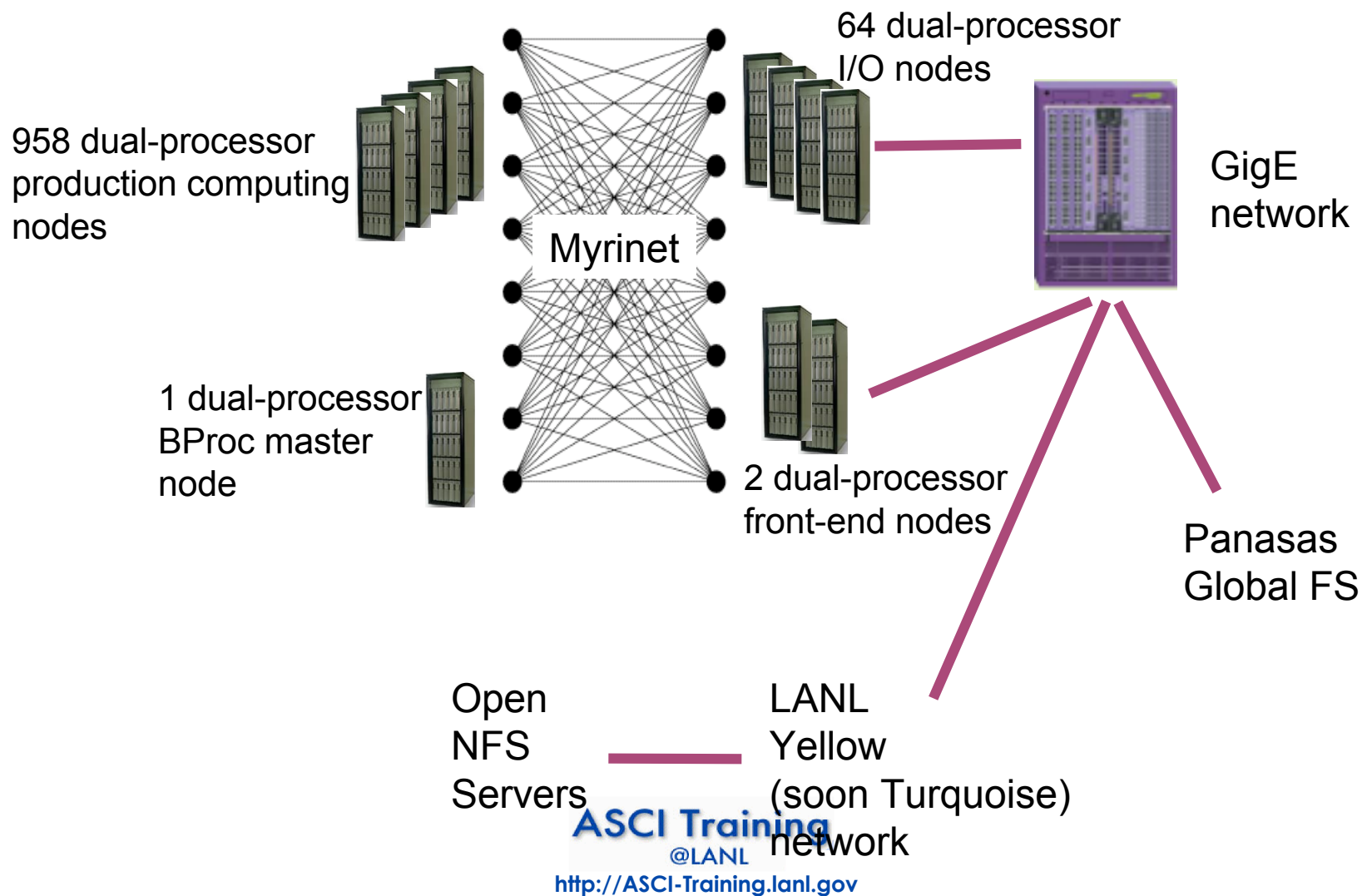


Pink Configuration

- **1,024 2-processor nodes**
 - some front-end nodes & 64 fileserver nodes.
- **Each node has:**
 - 2 Intel 2.4-GHz Xeon Processors / 2 GB Memory
 - zero disk drives, zero ethernet cables ***
 - *** except for front ends
 - DDR SDRAM (instead of RDRAM) & 400-MHz system bus
 - Peak memory BW 3.2 GB/s; 1.6 GB/s more typical (compare to 2.1 GB/s on Alpha 21264)
- **9.8-TeraOps peak system performance.**



Pink Configuration





Intel Xeon Microprocessor

- 32-bit x86 (IA32) architecture.
- Xeon is basically a multiprocessor version of Intel Pentium IV
- The x87 FPU uses a floating point stack with eight 80-bit register elements. These same 80-bit registers are used whether the operands are single, double, or extended-double precision.
- SIMD "vector" units: MMX & SSE 2
- `pink.lanl.gov-> more /proc/cpuinfo`
model name : Intel(R) Xeon(TM) CPU 2.40GHz
cache size : 512 KB



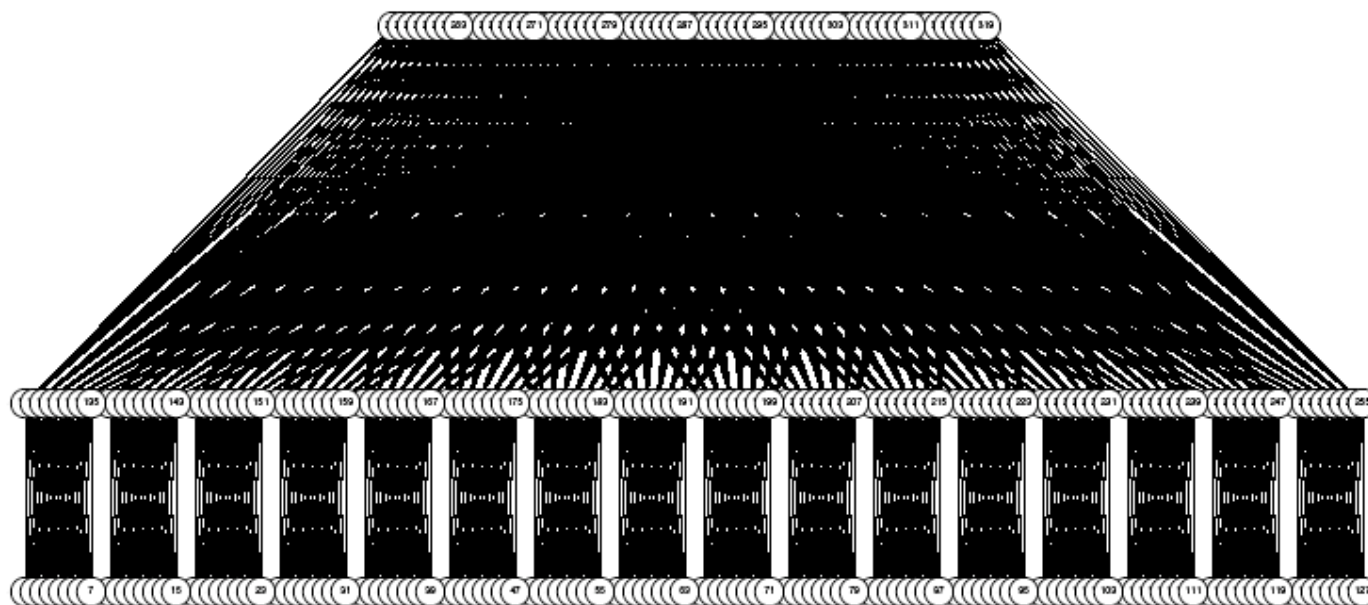
Myrinet Interconnect

- Network interface card connects to nodes' PCI I/O bus + series of federated 16-port cross-bar switches + optical cables + Myrinet gm software + CCS-1 mapper
- There is one "rail" of interconnect.
- Myrinet should have a unidirectional MPI send/receive latency of 6-7 microseconds and a peak unidirectional transfer bandwidth of 250 MB/s. See CCS-3 results for details.



Myrinet Interconnect

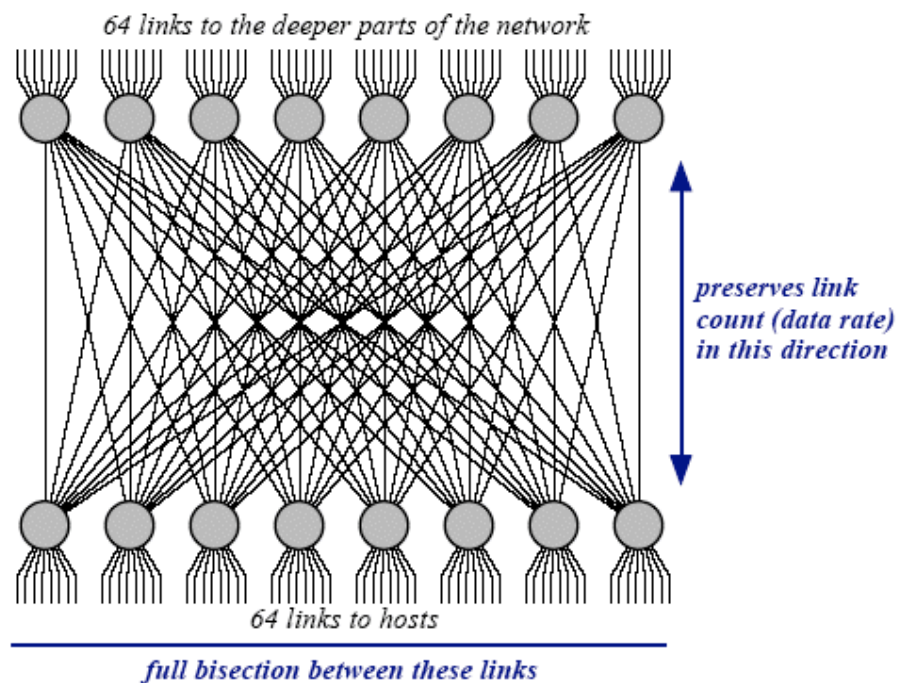
1024 Hosts, 320 SW16 Switches



- 5-hop worse-case transfer
- Bandwidth-preserving topology (similar to fat-tree)



Myrinet Interconnect





Pink Software Architecture

- Pink is an example of a "Science Appliance."
- Award-winning concept invented by LANL's CCS-1 Cluster Computing Team (Ron Minnich, TL).
- *Objective is to provide more computing cycles to users by making the cluster easier to build and manage.*
- Pink is essentially the world's largest Science Appliance. *
 - * Lightning has more nodes but is operated as 6 separate segments.

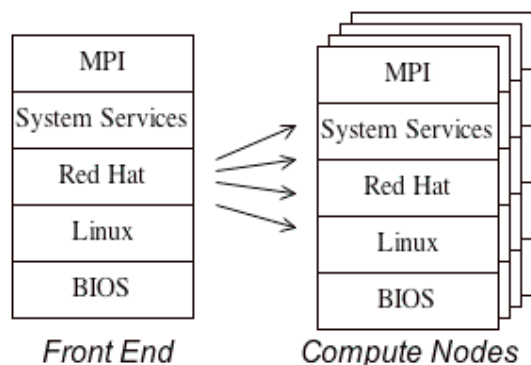


What is a Science Appliance?

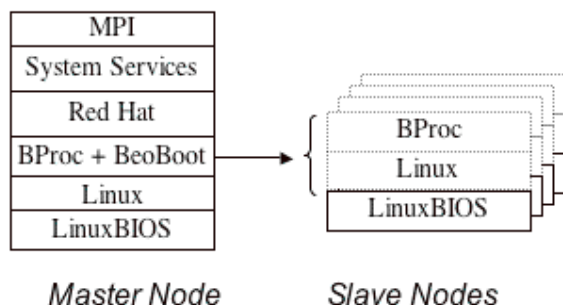
- Science Appliance actually refers to a redesign of both hardware and software for large-scale clusters.
- The key software in a Science Appliance is a suite that LANL developed called "Clustermatic."
 - Clustermatic can completely control a cluster, from the BIOS up to a high level programming environment.
 - It features the Beowulf Distributed Process Space (BProc), LinuxBios, and a variety of other open-source kernel modifications, utilities, and libraries.
- Reliability improved by reducing HW and SW complexity; availability improved by reducing boot time to seconds or a few minutes.



Science Appliance vs. a Traditional Cluster



Traditional Cluster
Architecture



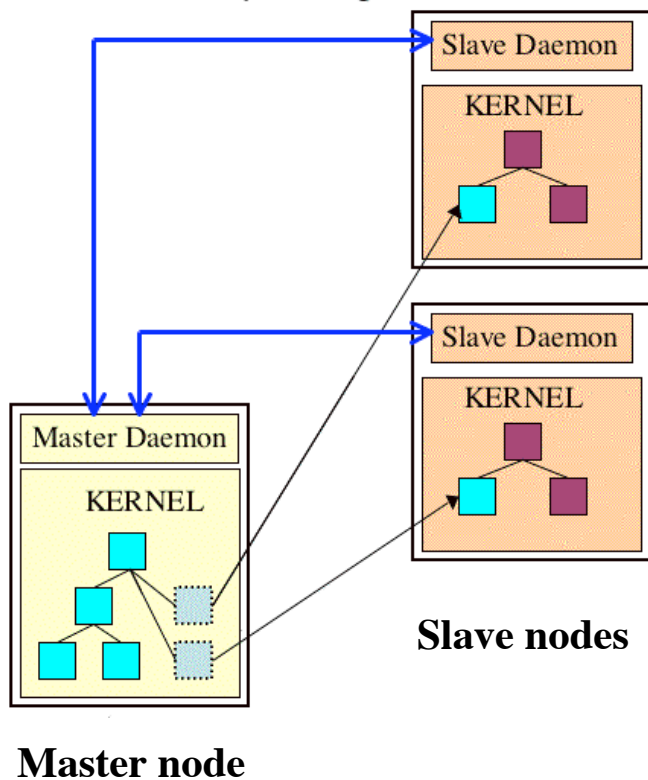
Science Appliance
Architecture

- A traditional cluster is built by replicating a complete workstation's software environment on every node.
- In a Science Appliance, we have master nodes and slave nodes but only the master nodes have a fully-configured system.
- The slave nodes run a minimal software stack consisting of LinuxBIOS, Linux, and BProc.
- No Unix shells running on the slave nodes, no user logins on the slave nodes.



Science Appliance vs. a Traditional Cluster

Process Tree Spanning 3 Machines

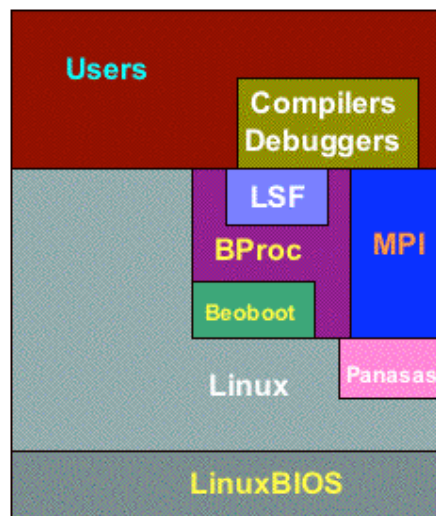


- Most importantly, BProc enables a distributed process space across nodes within the Pink cluster: all user processes running on the slave nodes appear as processes running on the front end.
- Users create processes on the master node and the system migrates them (the processes) to the slave nodes.
- Standard input, output, and error streams are redirected to the master node.

Processes remain visible, controllable on master.



Science Appliance Systems at LANL



CCS-1 Cluster Research Team
CCS-1 LA-MPI (Resilient Tech.) Team
3rd-Party Vendor

- Lightning, Pink, Grendels, Flash
- MPI & LSF are BProc-integrated.
- File I/O environment is different among the various LANL BProc systems
- Result: LANL Science Appliance systems are easy to use but are different than other LANL systems (hence this presentation).

- Important note: The root filesystem in Clustermatic is RAM based!



Porting Considerations

- Endianness
- Data Sizes
- 32-bit address space: 2-GB memory
- Large file support: Nominal limit is 2-GB files.
 - See "Big File Fix" .+3 view graphs



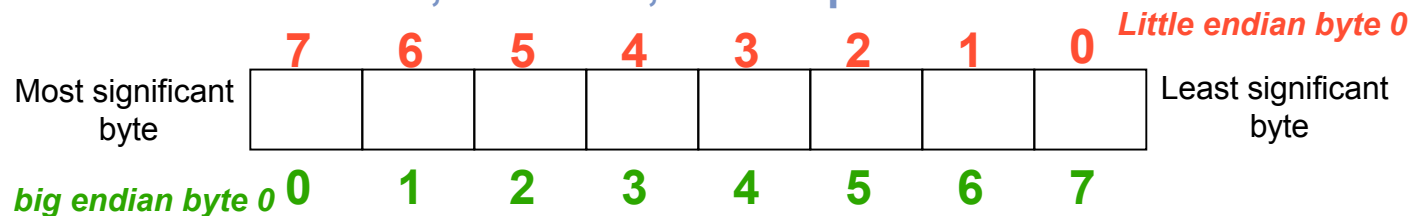
Porting Considerations: Endianness

Machine	Processor	Byte order
ASCI Q	DEC/Compaq/HP Alpha EV68	Little Endian
Lambda, Pink, Grendels	Intel Pentium	Little Endian
Blue Mountain or Theta	MIPS R10000	Big Endian
Mauve	Intel Itanium	Little Endian
Lightning	AMD Opteron	Little Endian



Porting Considerations: Endianness

- **Big Endian:** address of most significant byte = word address (big end of the word).
 - 0xDEADBEEF = DE AD BE EF
 - IBM 360, Power; Motorola 68k; MIPS; SPARC; HP PA
- **Little Endian:** address of least significant byte = word address (little end of the word)
 - 0xDEADBEEF = EF BE AD DE
 - Intel 80x86; DEC Vax; DEC Alpha



- **Code issues:** pointer dereferencing can work differently
- **Data file issues:** Data files written on Theta must be converted to little-Endian format before they can be read on Pink.
 - Some compilers have a compile-line option



Porting Considerations: Data Sizes

Fortran Data Type	Format	Range
INTEGER	2's complement integer	-2^{31} to $2^{31}-1$
INTEGER*2	2's complement integer	-32768 to 32767
REAL	Single-precision floating point	10^{-37} to 10^{38} (1)
REAL*4	Single-precision floating point	10^{-37} to 10^{38} (1)
REAL*8	Double-precision floating point	10^{-307} to 10^{308} (1)
DOUBLE PRECISION	Double-precision floating point	10^{-307} to 10^{308} (1)



Porting Considerations: Data Sizes

C and C++ Data Types (Size in Bytes)			
	Theta (-32)	Theta (-64)/QSC	Pink
short	2	2	2
int	4	4	4
long int	4	8	4
long	4	8	4
unsigned long	4	8	4
long long	8	8	8
double	8	8	8
float	4	4	4



Big File Fix

- The 32-bit address space limits files to 2 GB max.
- Linux kernel fixed this some time ago.
- Compile with
`-D_FILE_OFFSET_BITS=64 -D_LARGEFILE64_SOURCE=1`
`D_LARGEFILE_SOURCE=1`
- These preprocessor flags work with gcc and intel C as is.
- For PGI compilers add `-Mlfs`.
- Check that it worked with `'nm -B a.out'` and checking that things like `'open64'` are defined, not plain `'open'`.



Transferring Files

- HPSS via `k5psi` on `pfe1`, `pfe2`
- `psi` on `pink`, slave nodes
- `scp`
- However, plan is to move Pink to a new LANL network ("turquoise"); outside LANL firewall.



Modules

- The Modules package provides a convenient way for the system to make multiple versions of system software available and a convenient way for users to update their working environment.
- On Pink most system software - compilers, debuggers, parallel libraries - can be accessed only through the Modules package.
- The Modules package consists of two parts: a shell-level `module` utility/command and the modulefiles that the utility uses.
- On Pink you use the `module` utility on the front end.



Module Acts Differently on Pink

- Modulefiles have descriptive prefixes.
- You can use just the prefix in your commands to get the default version.

```
pink.lanl.gov-> module avail mpich  
mpich/1.2.5(default) mpich/1.2.5-intel mpich/1.2.5-pgi
```

```
pink.lanl.gov-> module load mpich
```

```
pink.lanl.gov-> module avail intel
```

```
intel/7.1(default) intel-c/8.1 intel-fortran/8.1
```

```
pink.lanl.gov-> module load intel
```

```
pink.lanl.gov-> module list
```

Currently Loaded Modulefiles:

1) mpich/1.2.5 2) intel/7.1



Pink Modulefiles (12/14/2004)

- Module avail

intel/7.1(default)

pgi/5.2-4

intel-fortran/8.1

pgi/5.1(default)

intel-c/8.1

totalview/6.4.0-2(default)

lamapi/1.5.6

lamapi/1.5.7

lamapi/1.5.8(default)

mpich/1.2.5(default)

mpich/1.2.5-pgi

mpich/1.2.5-intel

java2sdk/1.4.2_04



Compiling

- You will be compiling on the front-end system.
- Intel
 - Version 7.1 (default): one modulefile gets ifc (fortran compiler) and icc (C compiler)
 - Version 8.1: separate modulefile for ifort (fortran compiler) and icc
- Portland Group v5.1 and 5.2.4: pgf77, pgf90, pgcc
- Gnu: apparently don't need to load a modulefile



Intel Compilers

- `ifc` or `icc`, v7. `ifort`, v8.
- `-convert big_endian -convert little_endian`
- Optimization: `-O0`, `{0, -O1, -O2}` are equivalent on IA32, `-O3`
- `-tpp7`: Optimize build for Pentium 4 / Xeon.
- `-Vaxlib`: Fortran only - enables linking to compatibility library.
- `-g`: Build with debugging symbols. Intel lets you debug with optimization (`-O1` or `-O2`). `-O0` must be specified explicitly to turn all optimizations off.
- `-xW`: turns on auto-vectorizer (dubious value, IMHO)



Intel Compilers

- By default, all floating point exceptions (FPEs), such as invalid operation, denormal operand, overflow, underflow and divide-by-zero are masked on the IA32 architecture.
- Programs that encounter FPEs will not terminate unless steps are taken to unmask exceptions.
- No easy/direct way to catch FPEs with Intel `icc/ifc v7`. We are trying to import a method developed by LLNL to control FPEs. To use it, you will `#include` a header file `fpcontrol.h` and load with a library `libfpcontrol.a`. More information will be posted when available.



Compiling Miscellany

- **CPP is in /usr/bin/cpp**
- **Perl is in /usr/bin/perl (not where it is on Theta or QSC)**
- **Nice trick (supplied by Mike McKay, CCN-8):**
Instead of executing your perl scripts with a
hardwired directive line (e.g.,
#!/usr/bin/perl), use the following two lines as
the first two lines in your script:

```
eval 'exec perl -w -S $0 ${1+"$@"}'  
if 0;<rest of perl script>
```



MPI on Pink

- Two packages: LAMPI and MPICH
- **MUST** Load a modulefile, to compile **and/or run**
- Supports debugging with Totalview
- Must tell the compilers where to find include files and libraries if using LAMPI:

```
pfe1 -> module load lampi/1.5.6
pfe1 -> printenv |grep MPI
      MPI_ROOT=/usr/lampi-1.5.6/gm
      MPIHOME=/usr/lampi-1.5.6/gm
pfe1 -> pgf90 -I$MPI_ROOT/include \
        -L$MPI_ROOT/lib -lmpi *.o
```





MPICH on Pink

- Do not use the regular compilers (ifc/pgf90). Use the MPICH compiler scripts, instead (mpif90, mpicc).
- The modulefile automatically adds these to your path.
- **MUST** Load a modulefile, to compile **and/or run**
- Supports debugging with Totalview
- Additional runtime command line arg **REQUIRED**: `--nper 2`
- No need to add libraries to the link line.
- No need to tell the compilers where to find include files and libraries:

```
pfe1-> module load mpich/1.2.5-pgi
pfe1-> mpif90 *.o
pfe1-> mpirun -np 8 --nper 2 ./a.out
```



Additional MPI Note

- **Modulefiles have dependencies, i.e., mpich/1.2.5-pgi depends on pgi/5.2-4 or pgi/5.1(default)**
 - (Note difference in version #s)
- **Have to load both the MPI and compiler modulefile in order to run.**

```
pfe1 -> module load mpich/1.2.5-pgi
        mpich/1.2.5-pgi(19):ERROR:151: Module
        'mpich/1.2.5-pgi' depends on one of the module(s)
        'pgi/5.2-4 pgi/5.1 '
        mpich/1.2.5-pgi(19):ERROR:102: Tcl command
        execution failed: prereq pgi
pfe1 -> module load pgi mpich/1.2.5-pgi
```



Pink Filesystems

- Front end and back ends DO NOT have the same filesystems mounted.
- NFS Home directories (identical to lambda, QSC, etc.) mounted on front end ONLY.
- Front and back ends will have temporary workspace from Panasas filesystem:

`/net/scratch1/userid` and
`/net/scratch2/userid`

- Important: You want to do your runs in a directory that exists on both the front and back ends! Use the Panasas space.
 - May have to `cd /net/scratch1/user_id ; mkdir user_id`



Running Jobs on Pink

- Now the fun begins ...



Logging In

- Like most other LANL production computing systems, the Pink cluster uses a front end to control access to the compute nodes.
- You log in to the front end and then you submit jobs, both interactive and batch, to the back-end compute servers.
- You do not log in to the back ends.
 - But you still may need to use `llogin`.
 - No login sessions on BProc slave nodes.



WARNING: CHANGES COMING

- Access to the current front-end system, pink.lanl.gov, will be removed on December 22.
- Front ends will be pfe1.lanl.gov and pfe2.lanl.gov
- pink.lanl.gov will become a BProc master node only - user access only through LSF.
- This is to reduce extraneous load on BProc caused by compiling, editing, file xfer, etc.



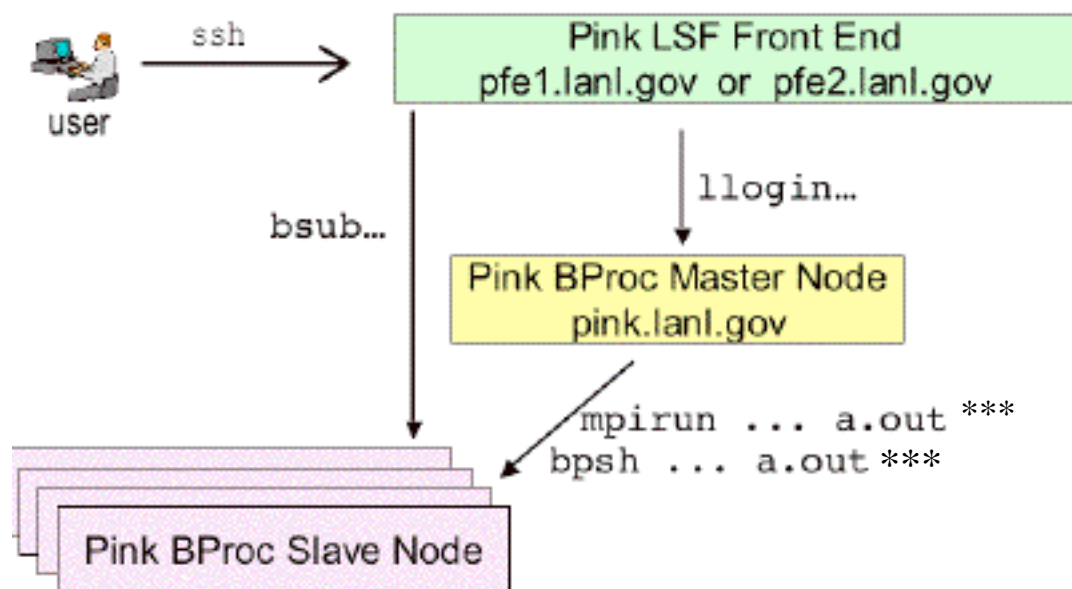
Running Jobs on Pink

- **LSF is on Pink and you must use it to run jobs.**
 - Both batch and batch-interactive jobs possible, as on other LANL systems.
- **If you don't use LSF your job runs on the front end!**
 - Different from other LANL systems.
- **LSF commands are basically the same as on other LANL systems.**
 - However, interactive use appears different.
- **The job submission process potentially involves a combination of LSF and BProc commands.**



Running Jobs on Pink

Procedure:



- Obtain an LSF allocation of slave nodes using `bsub` or `llogin`.
- Run job or script on master node.
- System migrates job to allocated slave nodes.

*** Illustration only; syntax will vary



Running Jobs on Pink

- **Terminology:**
 - "front ends" pfe1 and pfe2
 - LSF "hosts"
 - "submit_hosts" pfe1, pfe2, pink
 - "execution_hosts" pink
 - BProc "master node" (pink) and "slave nodes."
- There is one LSF execution host, comprised of all the BProc slave nodes available for computing.
- Use `llogin` to reach the BProc master node.



Obtaining a Slave Node Allocation with LSF

- On Pink you can never have a shell on a back-end system.
- If you choose to run interactively (`llogin` or `BSUB -Is`) you will be allocated slave node processors by LSF but your shell will still be on the BProc master node.

```
pfel.lanl.gov> llogin -n 8
Job <39681> is submitted to default queue <devq>.
<<Waiting for dispatch ...>>
<<Starting on pink>>
```

```
pink.lanl.gov> bjobs
```

JOBID	USER	STAT	QUEUE	FROM_HOST	EXEC_HOST	JOB_NAME	SUBMIT_TIME
39681	hjl	RUN	devq	pfel	8*pink	llogin	Dec 20 15:41



LSF on Pink

- Only the very basics considered here.
- If you don't know LSF, see <http://computing.lanl.gov> and documents referenced there.
- First we list the commands, then we talk about how to use LSF to get a slave node allocation.



LSF Commands: Host Status

- `lshosts` shows static resource information for the machines, such as number of CPUs, total memory, total swap space, etc.

```
pink.lanl.gov> lshosts
```

HOST_NAME	type	model	cpuf	ncpus	maxmem	maxswp	server	RESOURCES
pink	BPROC	PC2400	46.0	1892	2021M	-	Yes	()
pfe1	BPROC	PC2400	46.0	-	-	-	No	()
pfe2	BPROC	PC2400	46.0	-	-	-	No	()

- `lsload` gives dynamic load info



LSF Commands: Host Status

- `bhosts` shows mixed dynamic/static information about batch processing on the LSF host.

HOST_NAME	STATUS	JL/U	MAX	NJOBS	RUN	SSUSP	USUSP	RSV
pink	ok	-	1892	1426	1418	0	0	8

Max job slots per user

Max job slots per host

job slots started

jobs running

pink.lanl.gov> **bhosts -w**

Less than 1,916? Some nodes must be down



Other LSF Commands

- **bsub** **submits a job for batch or interactive execution**
Usage: `bsub [options] a.out`
`bsub [options] < bsub_scriptfile`
 - `-n # processors`. Only node-level alloc on Pink.
 - `-W [hours:]minutes` `-q queue_name`
 - `-e error_filename` `-o output_filename`
 - `-ls` **submits a batch interactive job**
- **llogin** **special version of bsub for interactive use.**
 - Puts you in \$HOME with clean environment
- **bhist -a** **shows history of jobs, including ones that finished.**
- **bpeek** **shows stdout & stderr for unfinished batch jobs.**
- **bkill** **sends a job the SIGINT and SIGTERM signals.**



Other LSF Commands

- **bqueues** displays information about LSF batch queues

Usage: bqueues

bqueues -l

bqueues -l queue_name

bqueues -u user_id

pink.lanl.gov-> **bqueues -u hjw**

NAME	PRIO	STATUS	MAX	JL/U	JL/P	JL/H	NJOBS	PEND	RUN	SUSP
devq	7	Open:Active	-	32	1	-	24	0	24	0
batchq	5	Open:Active	1856	1024	1	-	0	0	0	0

- **bqueues -l batchq**

DEFAULT LIMITS:

RUNLIMIT

360.0 min of pink

MAXIMUM LIMITS:

RUNLIMIT

1440.0 min of pink

PROCLIMIT

4 4 1024

ASCI Training

@LANL

<http://ASCI-Training.lanl.gov>



Obtaining a Slave Node Allocation with LSF

- Note: Although your shell will still be on the front end after using `llogin`, it will be a *new* shell.
- Load your modulefiles **after** using `llogin` (`module load module_file_name`).
- Recall that modulefile-related errors can be somewhat obscure. Example:

```
pfel -> module load lampi
```

```
pfel -> llogin
```

```
pink -> mpirun -np 8 sweep3d.mpi
```

Not enough nodes to allocate all processes

```
pink -> module list
```

No Modulefiles Currently Loaded.



Obtaining a Slave Node Allocation with LSF

- LSF sets two important environment variables when it gives you a slave node allocation:
- **NODES** specifies which nodes the LSF job can use.
- **NODELIST** lists the processors on each node that the LSF job can use.

```
pfel -> llogin -n 6
```

```
Job <39681> is submitted to default queue <devq>.
```

```
<<Waiting for dispatch ...>>
```

```
<<Starting on pink>>
```

```
pink -> bjobs
```

JOBID	USER	STAT	QUEUE	FROM_HOST	EXEC_HOST	JOB_NAME
39682	hjl	RUN	devq	pfel	6*pink	llogin

```
pink -> env |grep NODE
```

```
NODES=4,5,11
```

```
NODELIST=4,4,5,5,11,11
```



Pink Queue Structure

	Priority	CPUs per Job Def / Max	Runlimit Def/Max hours	Queue Job Limit	User Job Limit	Notes
devq	7	2/8	12/12	-	16	ONLY_ INTERACTIVE
smallq	6	2 / 256	4/12	768	256	
largeq	5	258 / 1600	6/12	1032	1600	
nightq	5	- / 1600	2 / 10	-	1600	RUN_ WINDOW: 20:00-8:00

USER_SHARES: [institution, 35000] [support, 15000] [others, 5000]



Running Jobs on Pink

- Now you know how to get an allocation and what that looks like.
- What commands do you use to run jobs?
- The job submission process potentially involves a combination of LSF and BProc commands:

	Interactive	Batch
Sequential job	<code>llogin</code> <code>bpsh node# a.out</code>	<code>bsub -n 2 'bpsh \$NODES a.out'</code>
Parallel job	<code>llogin -n #</code> <code>mpirun -np # a.out</code>	<code>bsub -n # mpirun -np # a.out</code>



Only 1 BProc Command You Need to Know

- **bps**
- Three others it couldn't hurt to know:

bpstat bpps bptop



The BProc bpsh Command

- Pronunciation: bee-pish
- Runs a command on a slave node. Used for sequential commands. (Use `mpirun` for mpi.)
- Usage: `bpsh [options] node# command [cmd args]`
- Example: Run `ls -al /tmp/hjw` on node 3
`bpsh 3 ls -al /tmp/hjw`
- Common usage:
`pink -> llogin`
`pink -> bpsh $NODES a.out < test_24.inp > test_24.out`
- Caution: Does not execute a shell on back end

```
pink -> bpsh $NODES cd /tmp/hjw
bpsh: cd: command not found
```



Running MPI Jobs on Pink

- **With LAMPI use**

```
mpirun -np # a.out.mpi
```

to launch the job (plus LSF).

- No BProc command needed.

- **Extra manager process per node**

- **With MPICH use**

```
mpirun -np # --nper 2 a.out.mpich
```

- **Caution: mpirun doesn't check for LSF allocation**

- Runs on front end if no allocation exists!
- Can oversubscribe allocation without telling you.



Determining Job/System Status on Pink

- **Unix commands:** `ps`, `top`
- **LSF commands:** `bjobs`
- **BProc commands:** `bpstat`, `bpps`, `bptop`
- **Big difference:**
 - LSF command works on front ends and on BProc master node.
 - Unix commands and BProc commands work only BProc master node.



Determining Job/System Status on Pink

- `ps -ef` or `ps -f -u hjw`: slave node processes have their command names surrounded by [square brackets]

UID	PID	PPID	C	STIME	TTY	TIME	CMD
hjl	30293	30093	1	09:13	ttyp7	00:00:00	mpirun -np 4 ./sweep3d.mpi
hjl	30294	30293	19	09:13	ttyp7	00:00:00	[sweep3d.mpi]
hjl	30295	30293	19	09:13	ttyp7	00:00:00	[sweep3d.mpi]
hjl	30298	30295	87	09:13	ttyp7	00:00:01	[sweep3d.mpi]
hjl	30299	30294	92	09:13	ttyp7	00:00:01	[sweep3d.mpi]
hjl	30300	30294	0	09:13	ttyp7	00:00:00	[sweep3d.mpi]
hjl	30301	30295	0	09:13	ttyp7	00:00:00	[sweep3d.mpi]

- Watch for MPI processes without square brackets mistakenly running on front end!
- Unix `top` command will NOT show back end processes with square brackets.
- Can use Unix `kill` command on front end to send signal to back end processes.



Determining Job/System Status on Pink

- **LSF** `bjobs` command

```
pfel.lanl.gov-> bjobs
```

JOBID	USER	STAT	QUEUE	FROM_HOST	EXEC_HOST	JOB_NAME	SUBMIT_TIME
356	hjl	RUN	devq	pfel	4*pink	llogin	Jun 17 09:04

```
pfel.lanl.gov-> bjobs -u all
```

JOBID	USER	STAT	QUEUE	FROM_HOST	EXEC_HOST	JOB_NAME	SUBMIT_TIME
354	daughto	RUN	devq	pfel	16*pink	island3	Jun 16 21:45
356	hjl	RUN	devq	pfel	4*pink	llogin	Jun 17 09:04

- Show jobs for a single user: `bjobs -u hjl`
- Show jobs that are pending and why: `bjobs -lp`



The BProc bpstat Command

- **BProc bpstat command:** Shows status of nodes
 - up node is up and available
 - down node is down or can't be contacted by master
 - boot node is coming up (running node_up)
 - error an error occurred while the node was booting

Node #s

```

Sample bpstat and bjobs Output

[hjw@pink]$ bpstat
0,2,12,13
251
1,3-11,14-53
27-34
56
57-58
59-255

down  ----- root    root
error  ---x----- root    root
up     ---x----- greene  desktop
up     ---x----- jnunez  desktop
up     ---x----- scb     desktop
up     ---x----- hjw     desktop
up     ---x----- root    root

[hjw@ll-2 ~]$ bjobs -u all
JOBID  USER   STAT  QUEUE  FROM_HOST  EXEC_HOST  JOB_NAME  SUBMIT_TIME
6503   gap    RUN   devq   pink       4*pink     *n/tcsh -1  Mar 15 15:11
6482   greene RUN   devq   pink       100*pink   llogin     Mar 15 13:29
6488   jnunez RUN   devq   pink       16*pink    llogin     Mar 15 14:16
6489   scb     RUN   devq   pink       2*pink     llogin     Mar 15 14:18
6503   hjw     RUN   devq   pink       4*pink     *n/tcsh -1  Mar 15 15:11
  
```

ASCI Training

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<http://ASCI-Training.lanl.gov>



Determining Job/System Status on Pink

- Special CCN-7 version of `bpstat` command:
`bpps`
- Only shows slave node activity.

`pfel -> bpps`

NODE	USER	PID	PGID	S	STIME	TIME	COMMAND
5	hjl	31009	31008	S	09:45	00:00:00	[sweep3d.mpi]
6	hjl	31010	31008	S	09:45	00:00:00	[sweep3d.mpi]
6	hjl	31013	31008	R	09:45	00:00:06	[sweep3d.mpi]
5	hjl	31014	31008	R	09:45	00:00:06	[sweep3d.mpi]
5	hjl	31015	31008	R	09:45	00:00:06	[sweep3d.mpi]
6	hjl	31016	31008	R	09:45	00:00:06	[sweep3d.mpi]
8	lpm	14822	14597	R	09:31	00:01:01	[mpihello]
8	lpm	14823	14868	R	09:31	00:01:01	[mpihello]

- Returns nothing if no slave node processes



Common BProc Pitfalls

- The BProc commands DO NOT work on the front-end systems pfe1 and pfe2.

```
pfe1.lanl.gov> bpstat  
bproc_nodelist: Input/output error
```

```
pfe1.lanl.gov> bpps  
Bproc::proclist: Cannot allocate memory
```

```
pfe1.lanl.gov> bpsch 3 ./sweep3d.single  
bproc_vexecmove_io: Function not  
implemented
```



Common BProc Pitfalls

- `bsub -n 2 bpsch $NODES a.out` **FAILS**

Result: `NODES: Undefined variable.`

Correct behavior is obtained by quoting:

```
bsub -n 2 'bpsch $NODES hostname'
```



Common BProc Pitfalls

- `bpsh 6 ls /tmp; ls/tmp/hjw` **FAILS** because the **semicolon** is a shell command separator.

Result: `ls /tmp` runs on node 6
but `ls /tmp/hjw` runs on the front end!

- `bpsh $NODES ls > ~hjw/out` Works, no problem, but immediately thereafter: `bpsh $NODES ls ~hjw/out` **FAILS**

Result:

`ls: /users/hjw/out: No such file or directory`

- **Common theme: shell interpretation of a bpsh command takes place on the FRONT END only.**



Common BProc Pitfalls

```
[hjl@pfe1 ~/SWEEP/PINK]$ llogin -n 6
Job <39887> is submitted to default queue <devq>.
<<Waiting for dispatch ...>>
<<Starting on pink>>
[hjl@pink ~]$ cd SWEEP/PINK
[hjl@pink ~/SWEEP/PINK]$ module load lampi
[hjl@pink ~/SWEEP/PINK]$ mpirun -np 6 ./sweep3d.mpi
LA-MPI: *** mpirun (1.5.9)
LA-MPI:client/adminMessage.cc:930:
adminMessage::serverConnect timeout 42 exceeded --
0 client sockets account for 0 processes!
LA-MPI was unable to start your application.
This may be because:
<snip>
```

No Slave Node Directory



Common BProc Pitfalls

No Slave Node Allocation

```
pfe1-> llogin
```

```
Job <39882> is submitted to default queue <devq>.
```

```
<<Waiting for dispatch ...>>
```

```
<<Starting on pink>>
```

```
pink-> bpsh 3 ./sweep3d.single
```

```
3: Operation not permitted
```

```
pink-> bpstat
```

Node(s)	Status	Mode	User	Group
157,278,701,740,778,889	down	-----	root	root
0-156,158-277,279-513	up	---x-----	mswarren	desktop
514	up	---x-----	esd	desktop
515	up	---x-----	hxv	desktop
516-517	up	---x-----	hjl	desktop



Debugging with Totalview

- Don't forget that both your source and executable must be visible on the slave nodes.
- The process on Pink differs for serial and parallel codes.



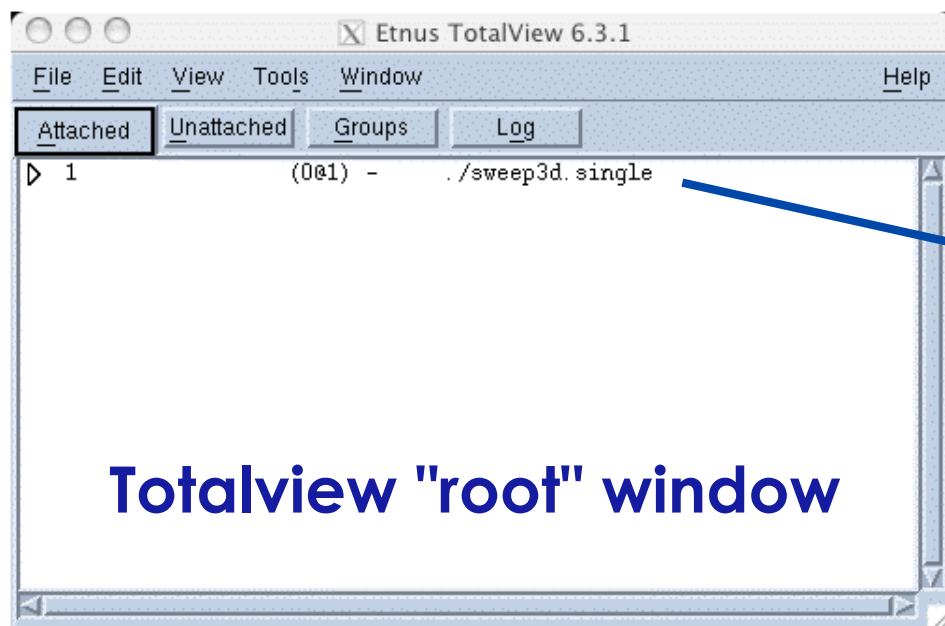
Debugging with Totalview on Pink

- Debug a serial code:

```
llogin
```

```
module load totalview/6.3.1
```

```
totalview -remote $NODES ./a.out
```



To start debugging,
double click on this line.

Totalview "root" window



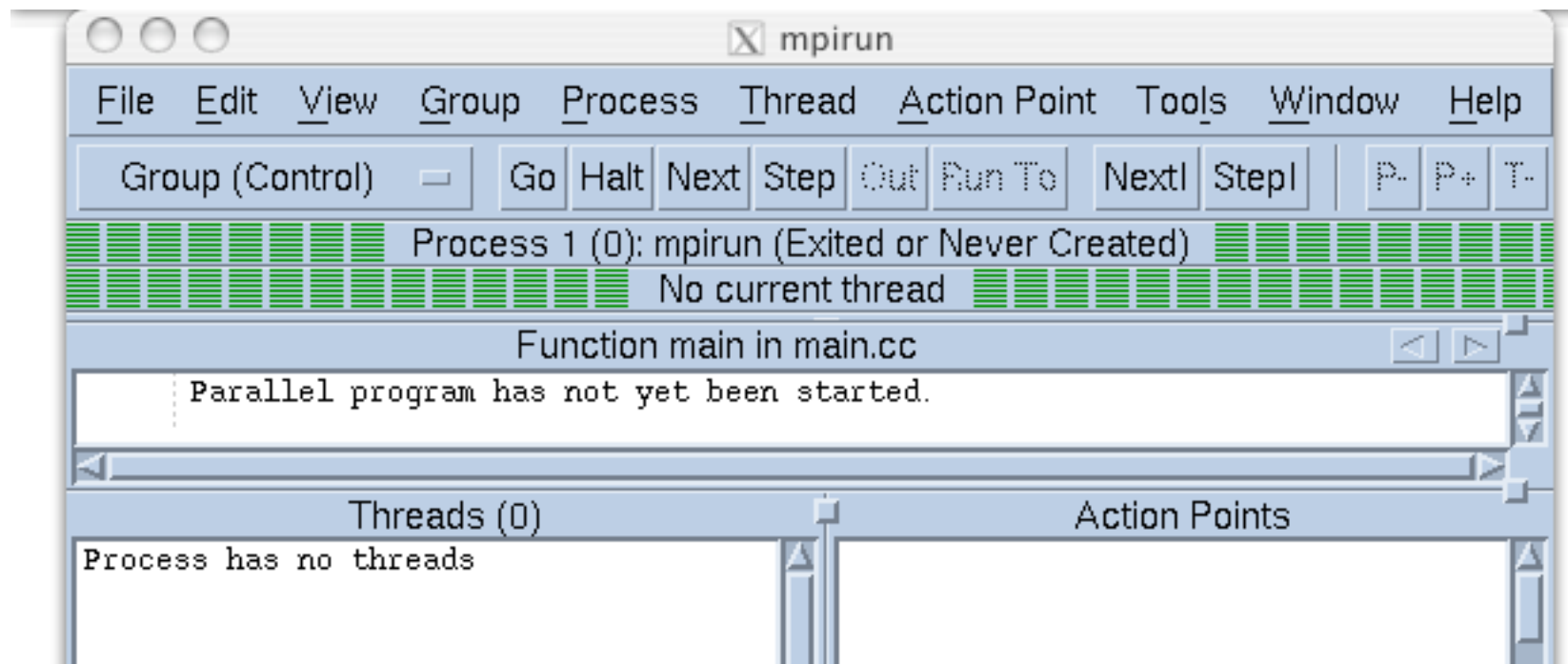
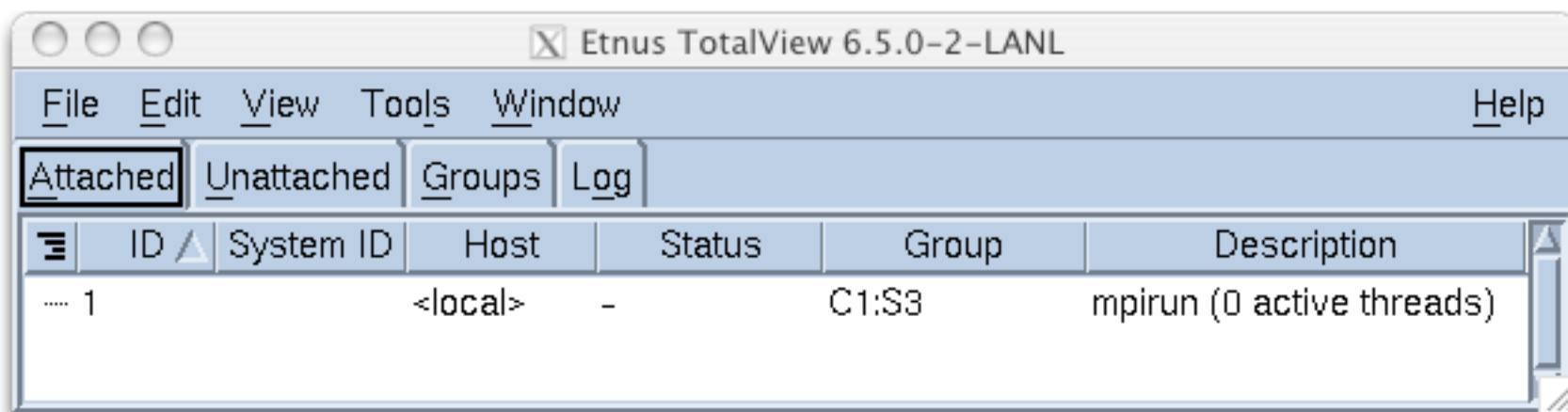
Debugging with Totalview on Pink

- Debug a parallel code:

```
llogin -n #  
module load lampi totalview  
totalview mpirun -a -np # ./a.out
```
- Currently a minor complication with LA-MPI: have to link in the additional binary "debug_gate.o" which is in the same directory pointed to by \$MPI_ROOT/lib.
- Debugging proceeds as it does on other LANL systems:



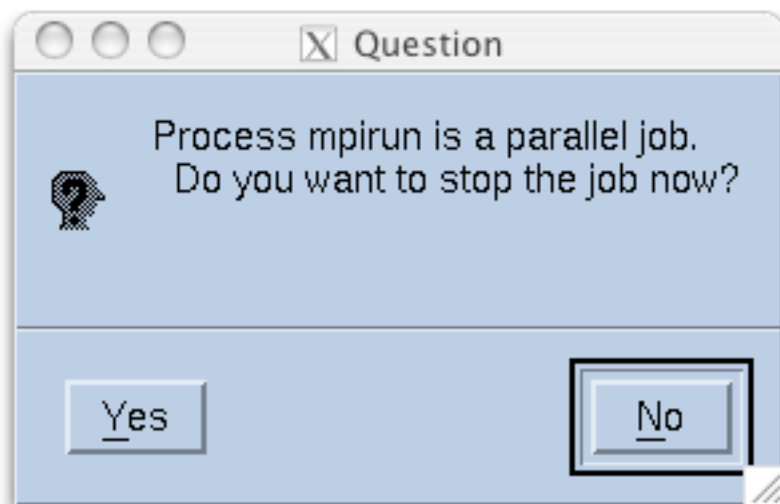
Totalview Root and Process Windows





Totalview Parallel Debugging

- Answer "Yes" to set breakpoints, see source, etc.





Panasas

- **Object Based Storage:** all I/O operations performed on arbitrarily-named data objects of variable size rather than sequentially numbered fixed-size blocks.
- The drives know about relationships among data objects. Reads and writes are directed to the object and an offset rather than to a logical block address.
- The Object Based Storage Devices stripe data across Secured storage devices. Rebalancing of the system is done automatically and is object based.
- To the user, the Panasas system should appear as a global shared file system with relatively high performance, even for small, sequential file I/O.
- Company founded by Garth Gibson.
<http://www.panasas.com>

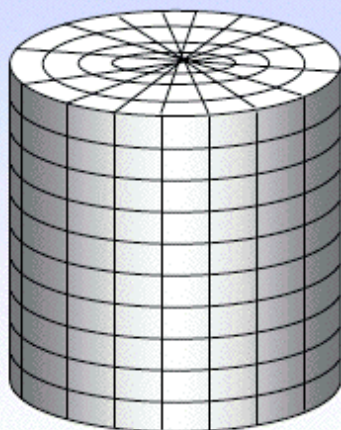


Panasas

Comparing Blocks To Objects

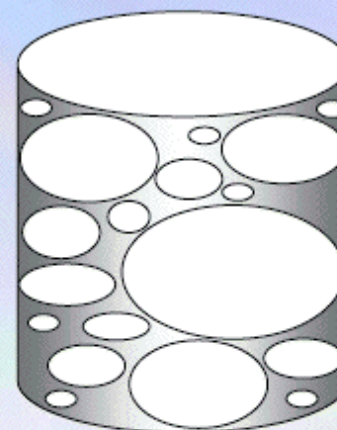
Block Based Disk

Operations
Read
Write
Addressing
LBA
Allocation
External

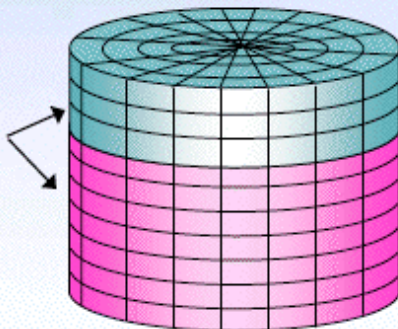


Object Based Disk

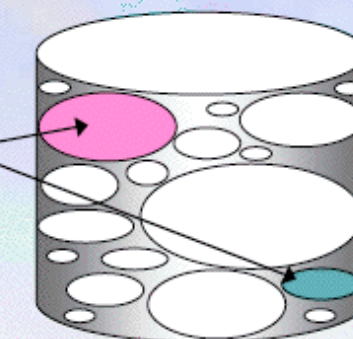
Operations
Create object
Delete object
Read object
Write object
Addressing
[Object, Offset]
Allocation
Internal



Security At
Volume Level



Security At
Object Level



ASCI Training

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Panasas

- In the Panasas system 10 Storage Blades and 1 Director Blade (containing the metadata managers) are combined into a single shelf containing 5 TB of storage and a 16-port Gigabit Ethernet switch (which uses 4 ports to the network and 11 to the blades). The blades use Intel Pentium processors and 250-GB ATA disk drives.



11-blade shelf
H: 7" [4U] x W: 19"

- On Pink, I/O goes over Myrinet to fileserver nodes, which then route I/O requests to Panasas. (Not that way on Lightning.) (Currently.)



Pink Summary

- `ssh pfe1.lanl.gov` **or** `ssh pfe2.lanl.gov`
- `llogin` **only if you need to run or query job status interactively (still have shell on front end)**
- `module load <compiler> <mpi>`
- `cd /net/scratch1/userid.` **(Run from here.)**
- `llogin + mpirun -np # a.out.mpi` **or**
`bsub -n # mpirun -np # a.out.mpi`
- `bpsh node# a.out.serial`
- **Make sure you don't execute a.out on front end!**
- `bpstat, bjobs, ps -ef, bpps`



Panasas User Best Practices

- See <http://computing.lanl.gov/article/439>
- Use `/usr/bin/panfs_df /net/scratch1` command instead of `df`.
- Small file performance is good, but highest bandwidths come from large I/O request sizes, larger than 100KB.
- For parallel N-to-1 I/O with you may be able to improve bandwidth using MPI / IO hint `mpio_concurrent_write` assuming your files do not overlap (or have ghost cells on the overlapping edges).



Where To Go For Help

- ICN Consultants support Pink users.
- 5-4444 option 3 or e-mail: consult@lanl.gov
- Do NOT send e-mail to sysadmins.
- <http://computing.lanl.gov>
- <http://asci-training.lanl.gov>
- Pink status eventually on
<http://icnn.lanl.gov>
and
<http://icnn.lanl.gov/drm/alljobs>



Want More Info?

- BProc Project Description <http://bproc.sourceforge.net>
- LANL's Cluster Research Team <http://public.lanl.gov/cluster>
- Clustermatic Home Page <http://www.clustermatic.org>
- LANL's Pink System <http://www.lanl.gov/projects/Pin>
- Linux Bios Home Page <http://www.linuxbios.org>
- A news article about Lightning from [fcw.com](http://www.fcw.com)
- Linux Networx's Home Page <http://www.linuxnetworx.com>
- Linux Networx's take on LinuxBios
linuxnetworx.com/products/linuxbios.php
- Scyld Computing Corp., a company selling BProc systems
- Linux Labs



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- Any errors in this presentation are due to hjw@lanl.gov only.